

Beyond the Edges of Kernel Control-Flow Hijacking Protection with HEK-CFI

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Contributions



- HEK-CFI: protection against kernel control-flow hijacking
- Proof-of-concept implementation
- Performance evaluation on Ubuntu 22.04 with an 1.85 % geomean overhead
- Security evaluation and comparison to other solutions

Motivation

Exploitation



Exploitation



Goals of adversaries

- Leaking sensitive informations, e.g., , , or 
- Resource compromising
- ...

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Kernel security

- Isolate different entities

Exploitation



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Kernel security

- Isolate different entities

Kernel vulnerabilities

- Exploitation to bypass isolation primitives

CVEs in the Linux Kernel

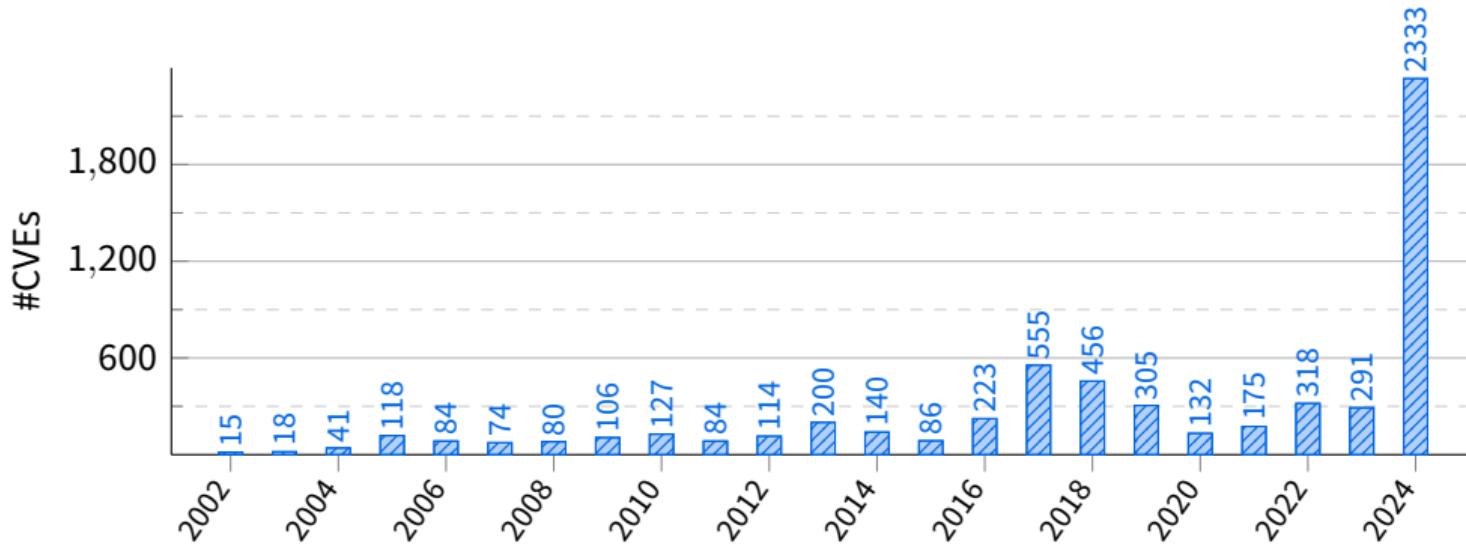


Figure: Found Linux kernel CVEs from NIST NVD.

Kernel Attacks



Kernel Attacks



Control-flow hijacking attacks

- Corrupt control data to redirect control flow
- Code execution → escalate privileges
- Popular, i.e., 15 out of 16 kernel exploits reported to Google's kernel bug bounty program [Goo22]

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Kernel Control-Flow Integrity (CFI) [CDA14, And22, ABEL05]

- Restricts the control flow to the Control-Flow Graph (CFG)
- E.g., Android ensures with function-signature granularity [And22]

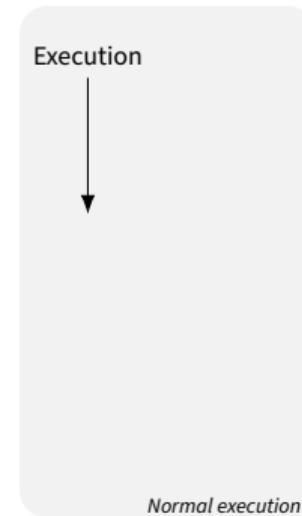
Motivational Example: CVE-2022-42703

- **CVE-2022-42703** [Set22] presents novel exploitation technique
 - Manipulates thread state for redirecting control flow



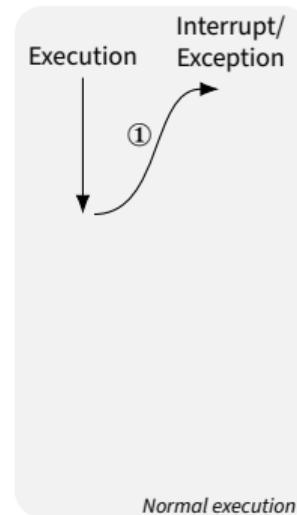
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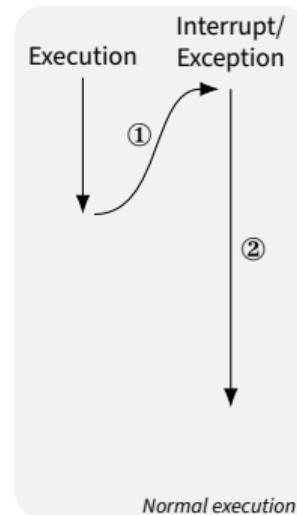
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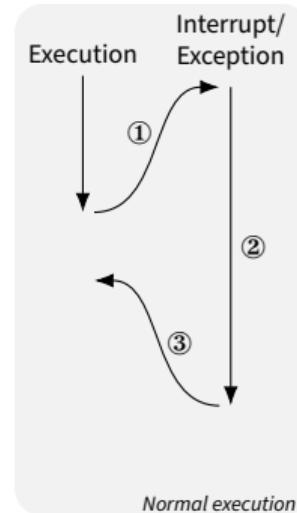
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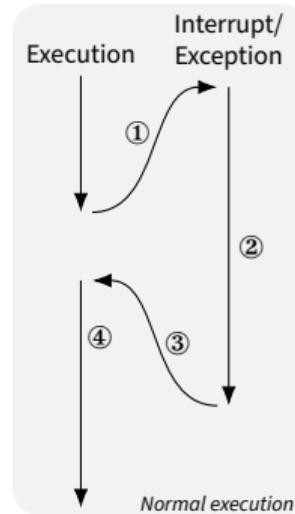
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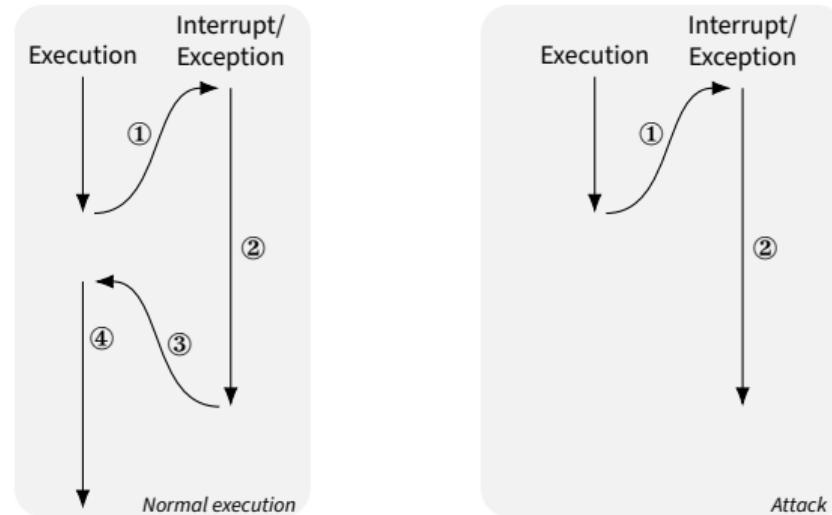
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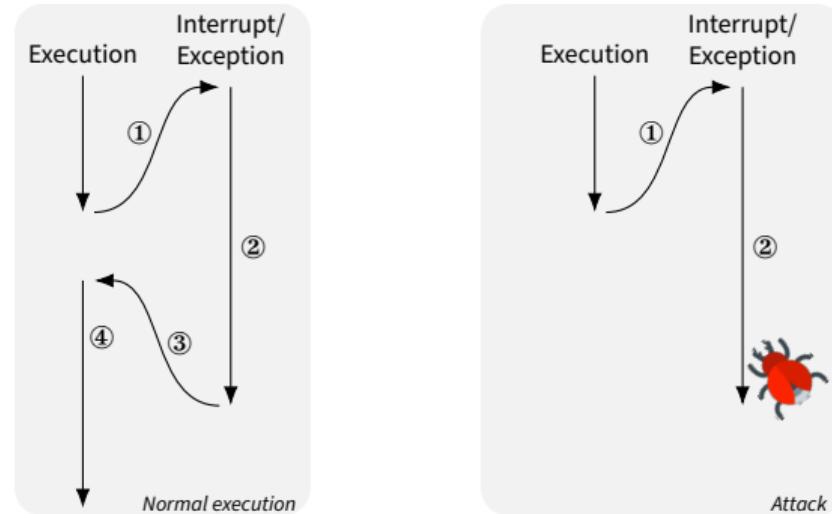
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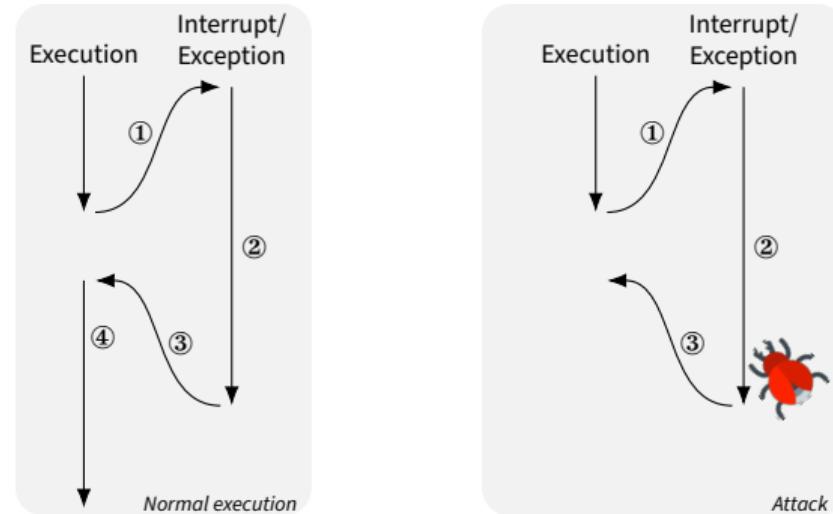
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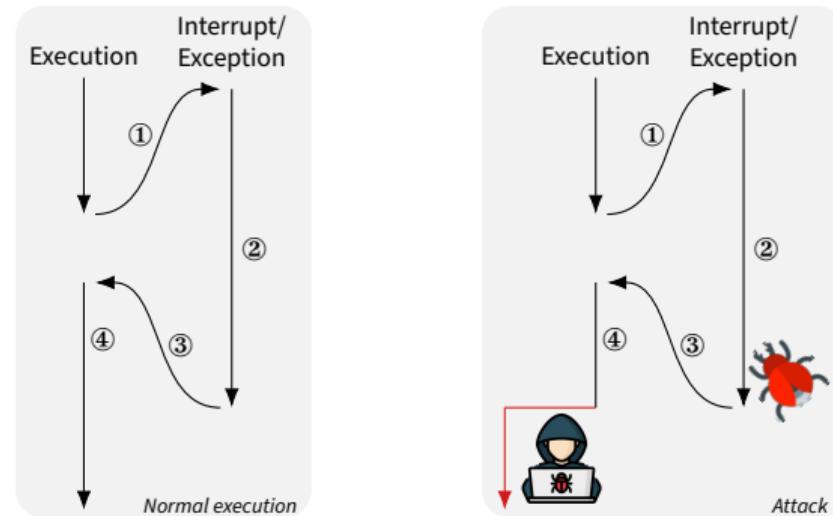
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Protecting Control-Flow Related Data

Control-Flow Hijacking Attacks in the Kernel

- ☒ Various control-flow related data allow to hijacking the control-flow

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```
1 struct timerfd_ctx {  
2     ...  
3     enum hrtimer_restart (*function)(struct hrtimer *);  
4     struct hrtimer_clock_base *base;  
5     ...  
6 }
```

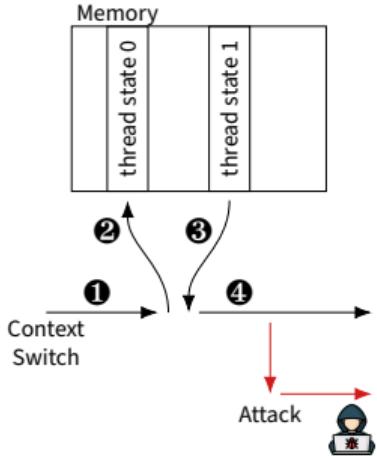
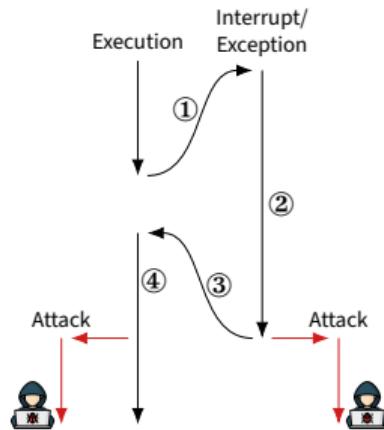
Control-Flow Hijacking Attacks in the Kernel

- ☒ Various control-flow related data allow to hijacking the control-flow
 - Function pointers
 - Operation table pointers

```
1 struct file_operations {  
2     ...  
3     ssize_t (*read)(struct file *, char *, size_t, loff_t *);  
4     ssize_t (*write)(struct file *, const char *, size_t,  
5                      loff_t *);  
6     ssize_t (*read_iter)(struct kiocb *, struct iov_iter *);  
7     ssize_t (*write_iter)(struct kiocb *, struct iov_iter *);  
8 };  
9  
10 struct file {  
11     ...  
12     const struct file_operations *f_op;  
13     ...  
14 };
```

Control-Flow Hijacking Attacks in the Kernel

- » Various control-flow related data allow to hijacking the control-flow
 - Function pointers
 - Operation table pointers
 - Thread state



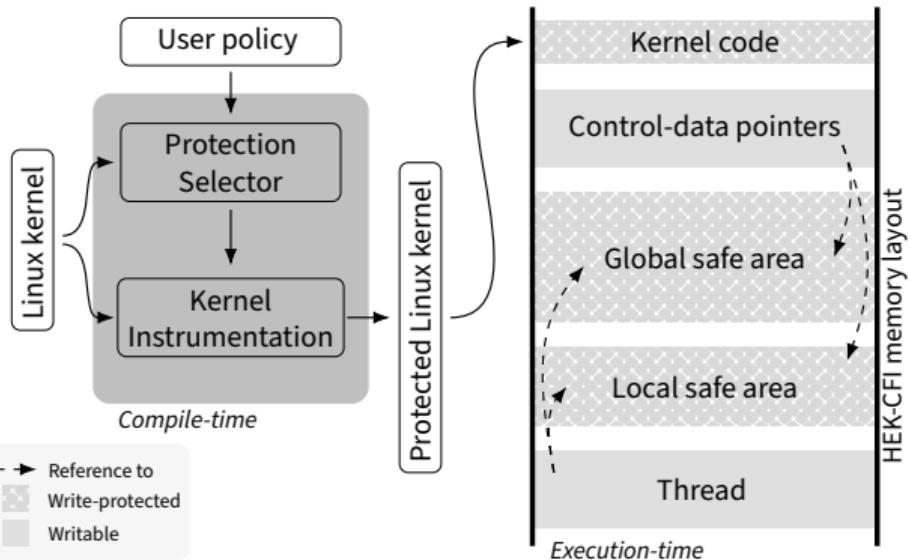
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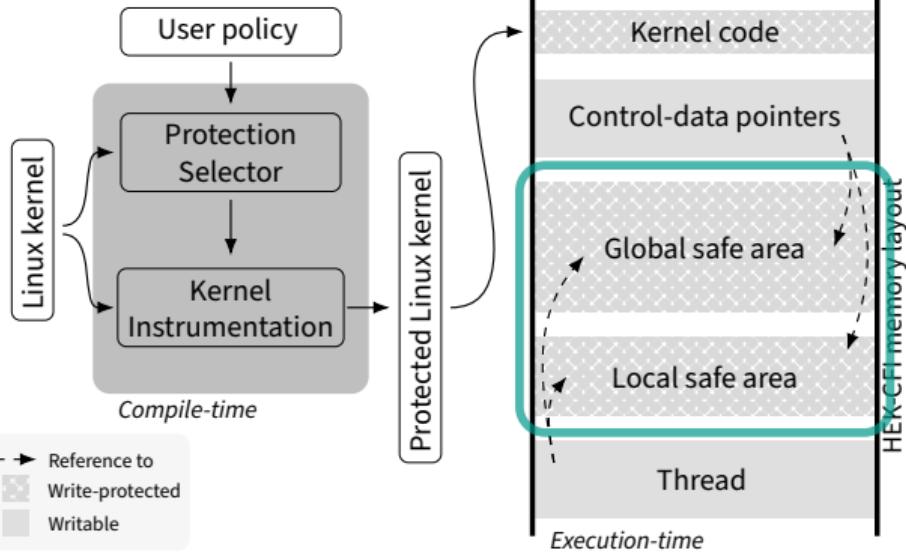
- Function pointers
- Operation table pointers
- Thread state
- Return addresses

```
1  dummy_fn:  
2  push  %r14  
3  push  %r13  
4  push  %r12  
5  mov   %rdi,%r12  
6  push  %rbp  
7  sub   $0x8,%rsp  
8  ...  
9  add   $0x8,%rsp  
10 pop   %rbp  
11 pop   %r12  
12 pop   %r13  
13 pop   %r14  
14 ret
```

Hardware Enforced Kernel Control-Flow Integrity (HEK-CFI)

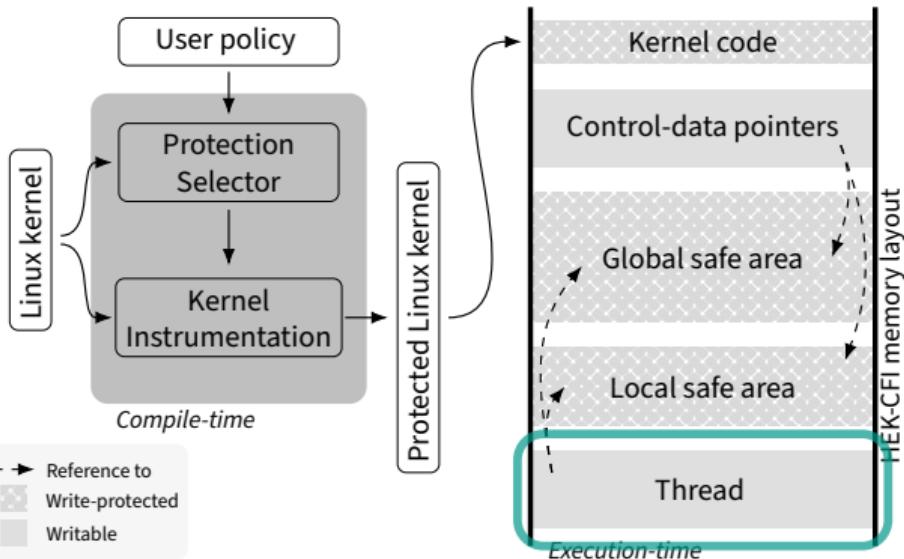


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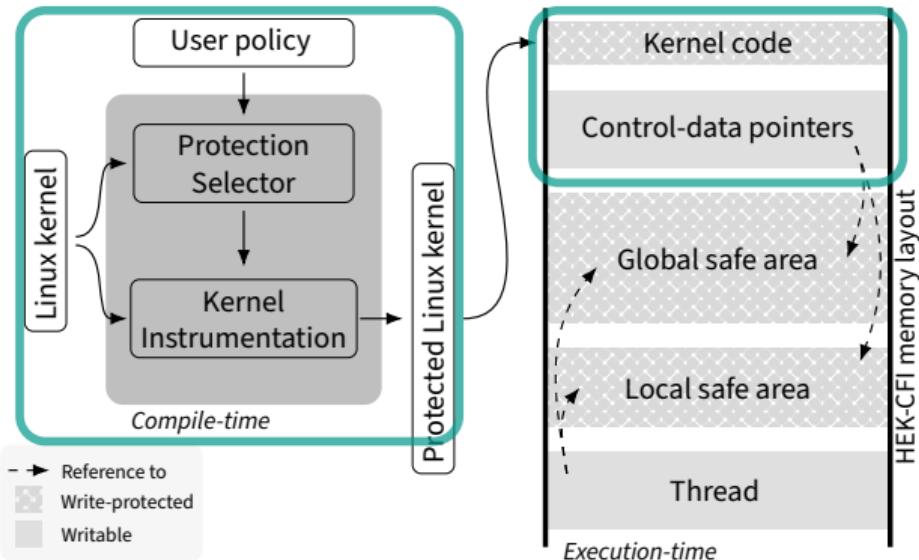
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 - Provides global/local safe areas
 - Write-protected memory areas

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 - Write-protected memory areas
- **Thread state protection**
 - Protects thread state with control-data integrity
 - Protects return addresses

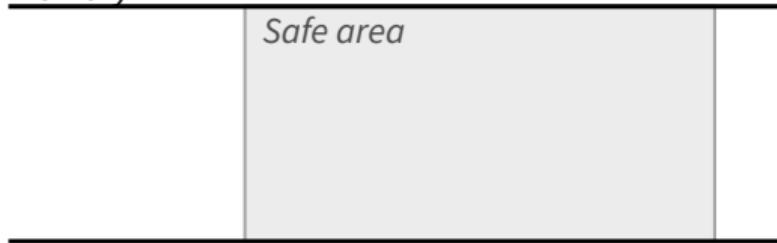
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- **Kernel control-data integrity**
 - Provides global/local safe areas
 - Write-protected memory areas
- **Thread state protection**
 - Protects thread state with control-data integrity
 - Protects return addresses
- **Protection selector & instrumentation**
 - Protects valuable pointers with control-data integrity
 - Protects non-valuable pointers with signature-based CFI
 - Based on user policy

Write-Protected Memory

Memory



Write-Protected Memory

Memory



- Write-protected memory for safe areas

Write-Protected Memory

Memory

	<i>Safe area</i>	
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 - Only certain instructions allowed to write to shadow pages [XWZ⁺22]
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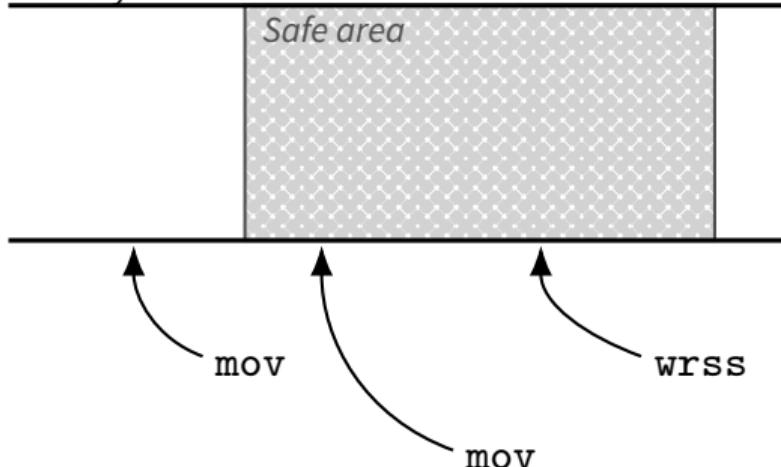


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 *write-protected by marking as shadow pages*

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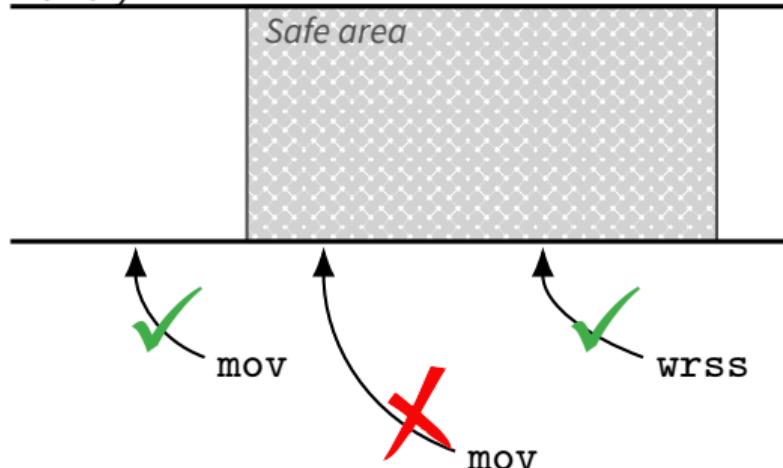


 write-protected by marking as shadow pages
→ write operation

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 - wrss for legal writes
 - Write operations (e.g., mov) causes exceptions

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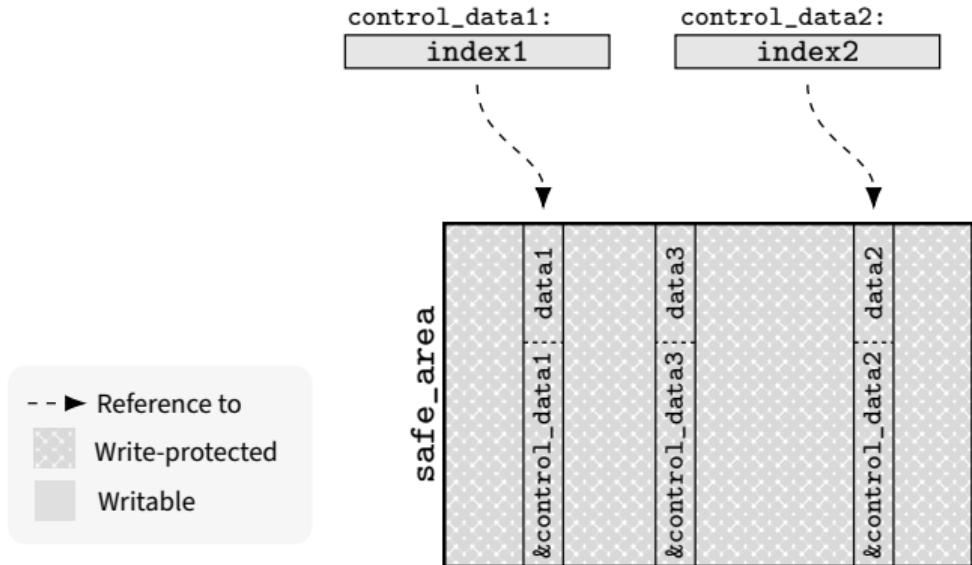
Memory



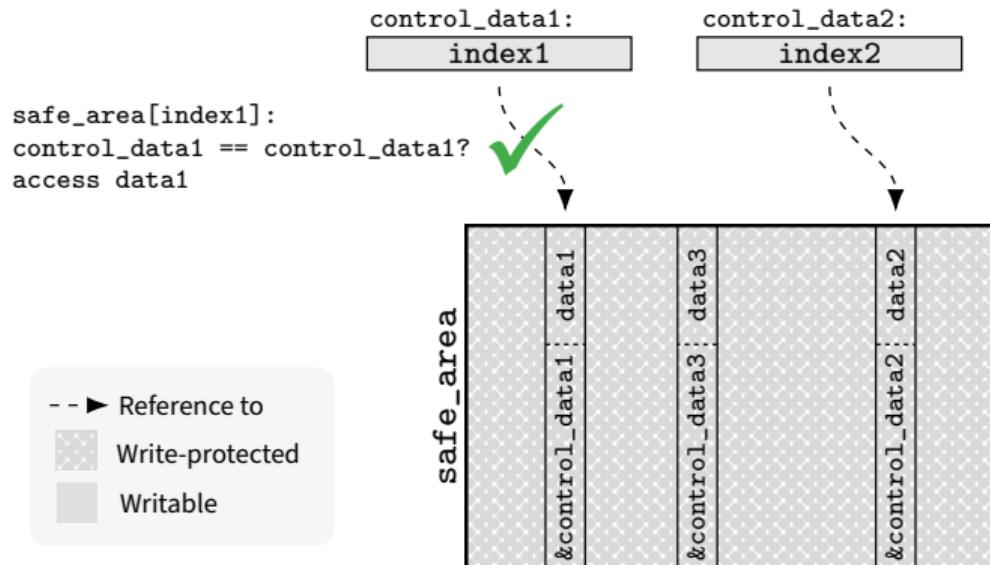
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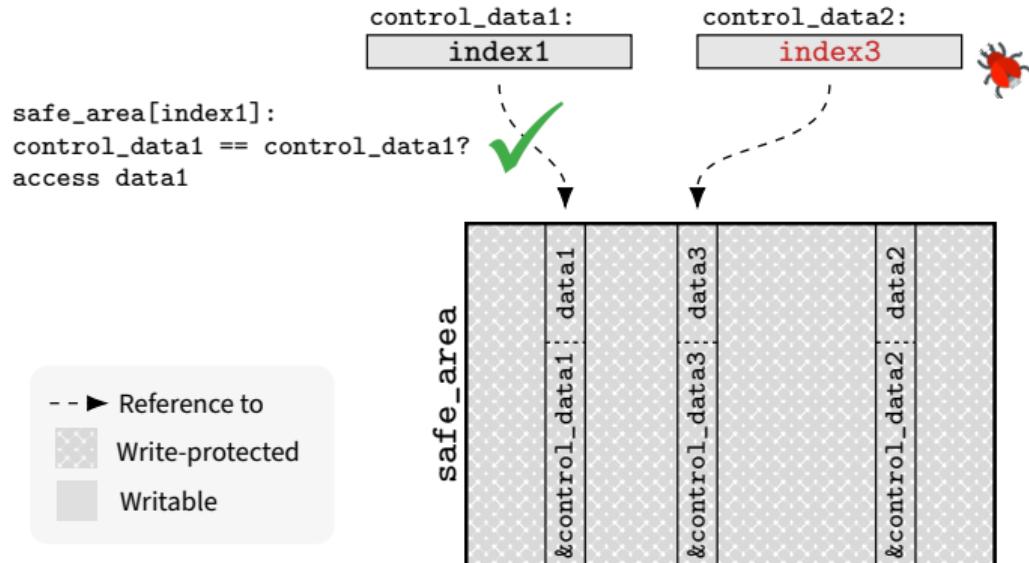
Global Safe Area



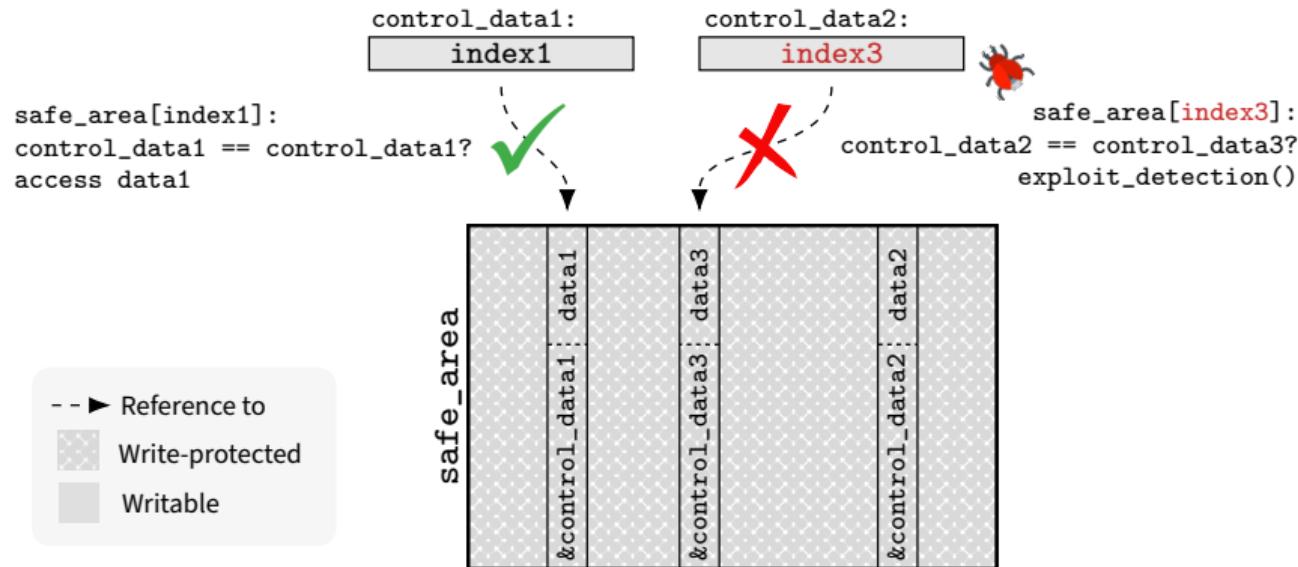
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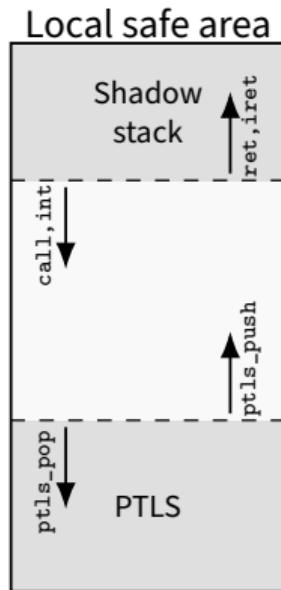
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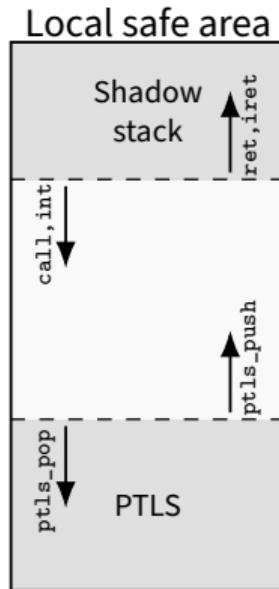
Global Safe Area



Local Safe Area

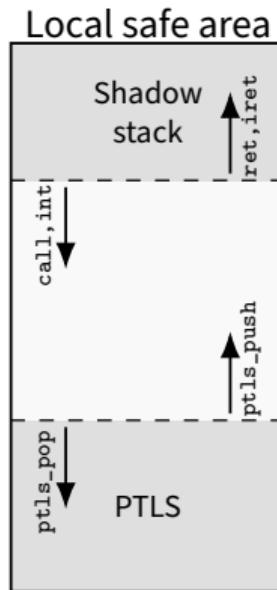


Local Safe Area



- Intel CET SHSTK
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Local Safe Area



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 - Use ss instructions to write to shadow stack
 - E.g., call pushes return addr
 - Does not provide pushss or pullss
- Protected Thread Local Storage (PTLS)
 - Software solution using wrss
 - Provides ptls_push/pull
 - E.g., used to stored thread state during interrupts/exceptions

Implementation & Evaluation

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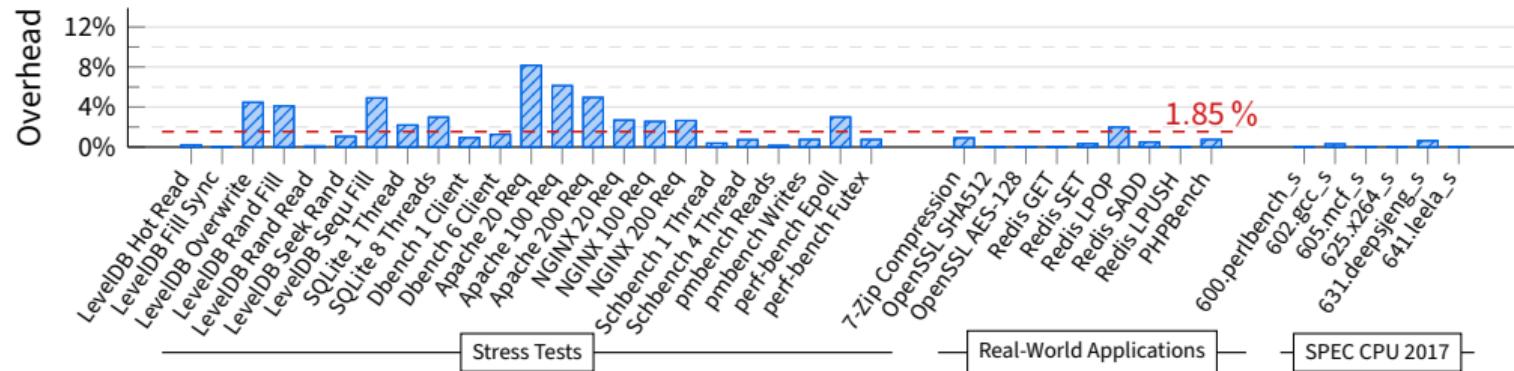
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Comparison

Mitigations	Attack Vector			
	Thread state	Return addresses	Operation table pointers	Function pointers
Ge et al. [GTPJ16]	○	○	■	●
kCFI [And22]	○	○	□	○
Fine-CFI [LTZM18]	○	○	■	●
PATTER [YZS ⁺ 19]	○	●	○	●
Camouflage [DCLCE20]	○	●	●	○
PAL [SJS ⁺ 22]	○	●	○	●
FineIBT [Mor22]	○	○	□	○
KCoFI [CDA14]	●	○	□	○
Intel CET SHSTK [Int16]	○	●	○	○
CPI [KSP ⁺ 14] + CETIS [XWZ ⁺ 22]	○	●	○	●
HEK-CFI	●	●	●	●

● Protection ○ Insufficient protection
 ■ Implicit protection □ Implicit insufficient protection
 ○ Does not protect but can be extended.

Conclusion



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- HEK-CFI that combines our kernel control-data integrity with signature-based CFI
- Implemented Intel CET SHSTK and a HEK-CFI proof-of-concept
- Performed a security and performance evaluation of HEK-CFI.

References I

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